

Resumen de Hopfield

Definimos una red con N neuronas binarias tipo Ising ($+1$ o -1), cuya dinámica está dada por una regla inspirada en las neuronas de McCulloch y Pitts pero **sin umbrales**.

$$S_i(t + \Delta t) = \text{signo}(h_i(t))$$

$$= \text{signo}\left(\sum_{j \neq i}^N w_{ij} S_j(t)\right)$$



dinámica secuencial $\Delta t = \frac{1}{N}$

" " paralela $\Delta t = 1$

De entre las 2^N posibles configuraciones

$$\bar{S} = (S_1, \dots, S_N)$$

de la red, nosotros escogemos el azar
con equiprobabilidad e independencia

p configuraciones que deseamos que
sean "estructuras" de la dinámica

$$\bar{S}^M = (\sum_{i=1}^N S_{i1}, \sum_{i=2}^N S_{i2}, \dots, \sum_{i=N}^N S_{iN}) \quad M=1 \dots p$$

Regla de Hopfield

$$W_{ij} = \begin{cases} \frac{1}{N} \sum_M S_i^M S_j^M & i \neq j \\ 0 & i = j \end{cases}$$

Duda bien si $\frac{P}{N} \approx 0.138$

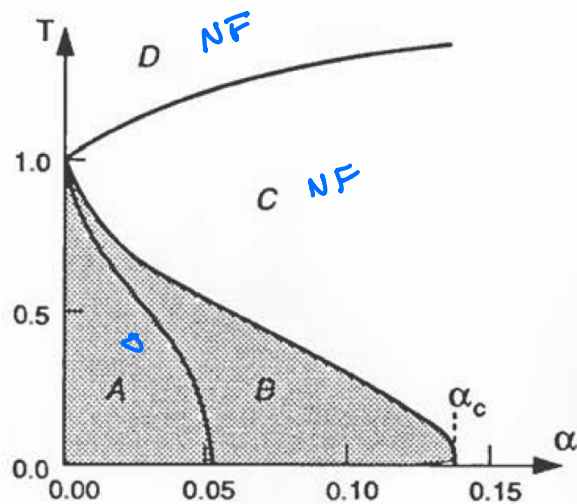
Podemos usar una dinámica estocástica

$$S_i(t+\Delta t) = \begin{cases} +1 & \text{con prob. } g(h_i) \\ -1 & \text{con prob. } (1-g(h_i)) \end{cases}$$

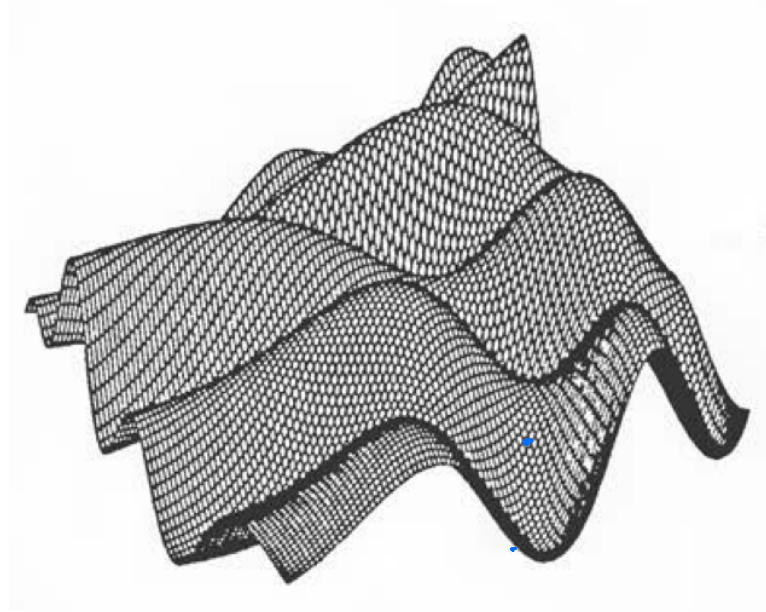
con

$$g(h) = \frac{1}{1 + e^{-\frac{2h}{T}}}$$

$T=0$ determinístico
 $T \rightarrow \infty$ $S = +1$ y -1 con $p = \frac{1}{2}$



$\approx \frac{0.12}{2}$



Per cada link $e_{ij}(t) = \frac{w_{ij} S_i(t) S_j(t)}$

$$\begin{aligned} E(t) &= \sum_{\text{links}} e_{ij}(t) \\ &= \sum_{i=1}^N \left[\sum_{\substack{j=1 \\ j \neq i}}^N w_{ij} S_i(t) S_j(t) \right] \end{aligned}$$

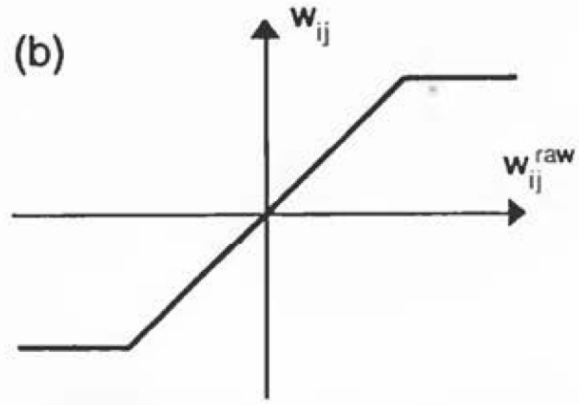
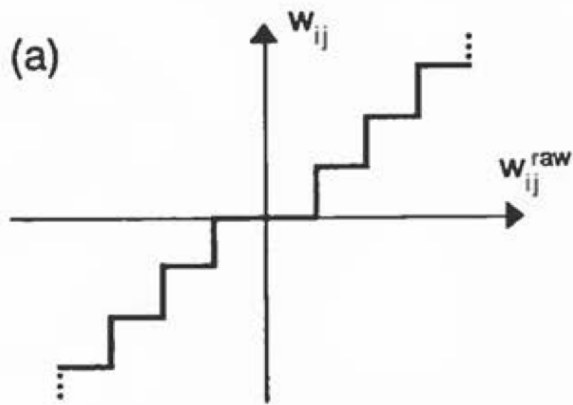
$$E(t + \Delta) \leq E(t) \quad T = 0$$

$$E = \sum_{i=1}^n \left[\sum_{\substack{j=1 \\ j \neq i}}^n \omega_{ij} S_j(t) \right] S_i(t)$$

$$E = \sum_{i=1}^n h_i(t) S_i(t)$$

Variaciones del modelo de Hopfield

o Clipping



o Aprendizaje

$$w_{ij}^{\text{nuevo}} = w_{ij}^{\text{antigos}} + \gamma \sum_i^{\mu} \sum_j^{\mu} \mu = 1, \dots, P$$

o Redes recurrentes ($\Delta t = 1$)

o Dilución débil

$$W_{ij} = \begin{cases} \frac{1}{N} \sum_{\mu} \sum_{i=1}^{\mu} \sum_{j=1}^{\mu} & \text{prob} = c \\ 0 & \text{prob} = 1 - c \end{cases}$$

$$C_{ij} = C_{ji} \quad \text{o} \quad C_{ij} \neq C_{ji}$$

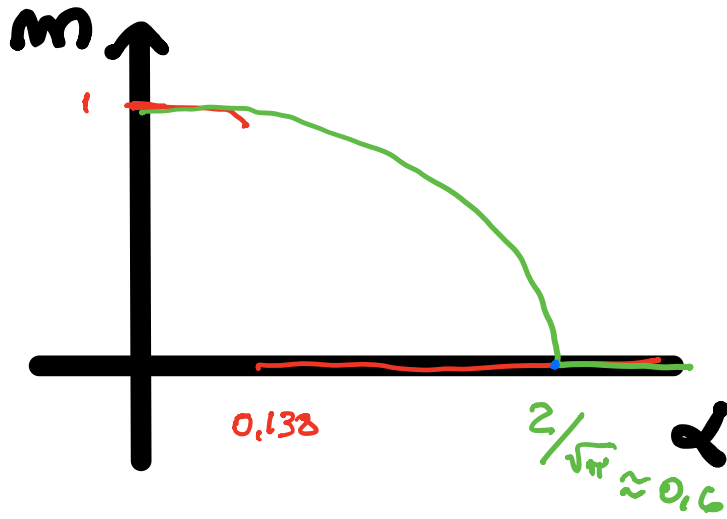
o Dilución fuerte

$$W_{ij} = \frac{1}{K} C_{ij} \sum_{\mu} \sum_{i=1}^{\mu} \sum_{j=1}^{\mu}$$

David
Gardner
Zipfelius

$$K \ll \ln(N)$$

$$C_{ij} = \begin{cases} 1 & \text{con prob. } \frac{K}{N} \\ 0 & \text{" " } \left(1 - \frac{K}{N}\right) \end{cases}$$

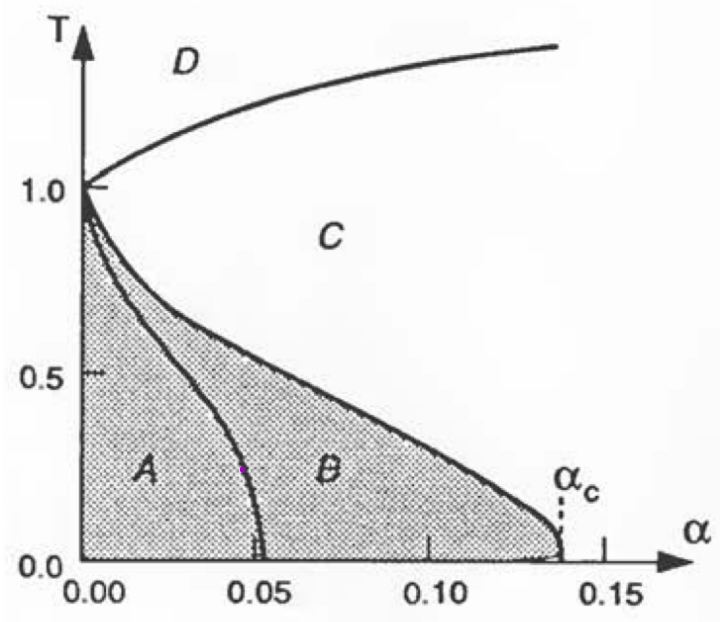


$$m = \frac{1}{N} \sum_{i=1}^N \underbrace{s(x_i)}_{0,4}$$

$$m = \operatorname{erf} \left(\frac{m}{\sqrt{2\alpha}} \right)$$

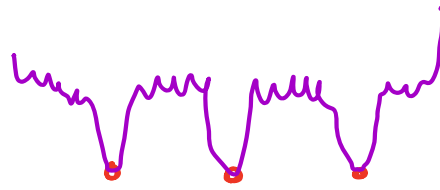
$$\alpha = \frac{k}{2}$$

$$\operatorname{erf}(x) = \frac{2}{\sqrt{\pi}} \int_0^x e^{-t^2} dt$$

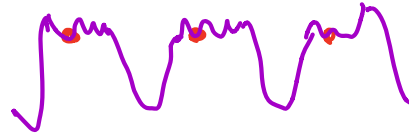


$\underline{P} \rightarrow 0$

N



A



B

$$F(T, N) \propto - \ln (Z(N, T))$$

$$\lim_{N \rightarrow \infty} \frac{F(T, N)}{N} = f(T) = - \frac{1}{N} \ln (Z(T))$$

$$Z_{\{\xi\}}(T)$$

$$\langle\langle f \rangle\rangle = \frac{1}{N} \langle\langle \ln Z_{\{\xi\}}(T) \rangle\rangle_{\{\xi\}}$$

↑

trick de réplique Replica trick

$$\ln(x) = \lim_{n \rightarrow 0} \left(\frac{1 - x^n}{n} \right) \quad \text{à ?}$$